

# Waters signatures

1. Waters signatures... (choose as many options as you think apply)

- (A) ...are unique (in the sense that for every pk and message  $m$ , there is at most one valid signature  $\sigma$ )
- (B) ...are randomized (in the sense that the signing algorithm uses random coins to choose certain values)
- (C) ...are EUF-CMA secure in the standard model (i.e., without random oracles) under the CDH assumption

2. Recall that Waters signatures are of the form  $\sigma = (g^r, g^\alpha H(m)^r)$  for the secret key  $\alpha$  and a random  $r$ . How can a signature be re-randomized? That is: given  $(g^r, g^\alpha H(m)^r)$ , how can we generate another signature  $(g^{r'}, g^\alpha H(m)^{r'})$  for  $m$  for a fresh, uniformly distributed  $r'$ , without knowing  $\alpha$ ?

Set  $r' = r + s$  for a known  $s$ . Then  $g^{r'} = g^r g^s$  and  $g^\alpha H(m)^{r'} = g^\alpha H(m)^r H(m)^s$  can both be computed without knowing  $\alpha$ .